Memory Forensics Current Practices and Future Directions

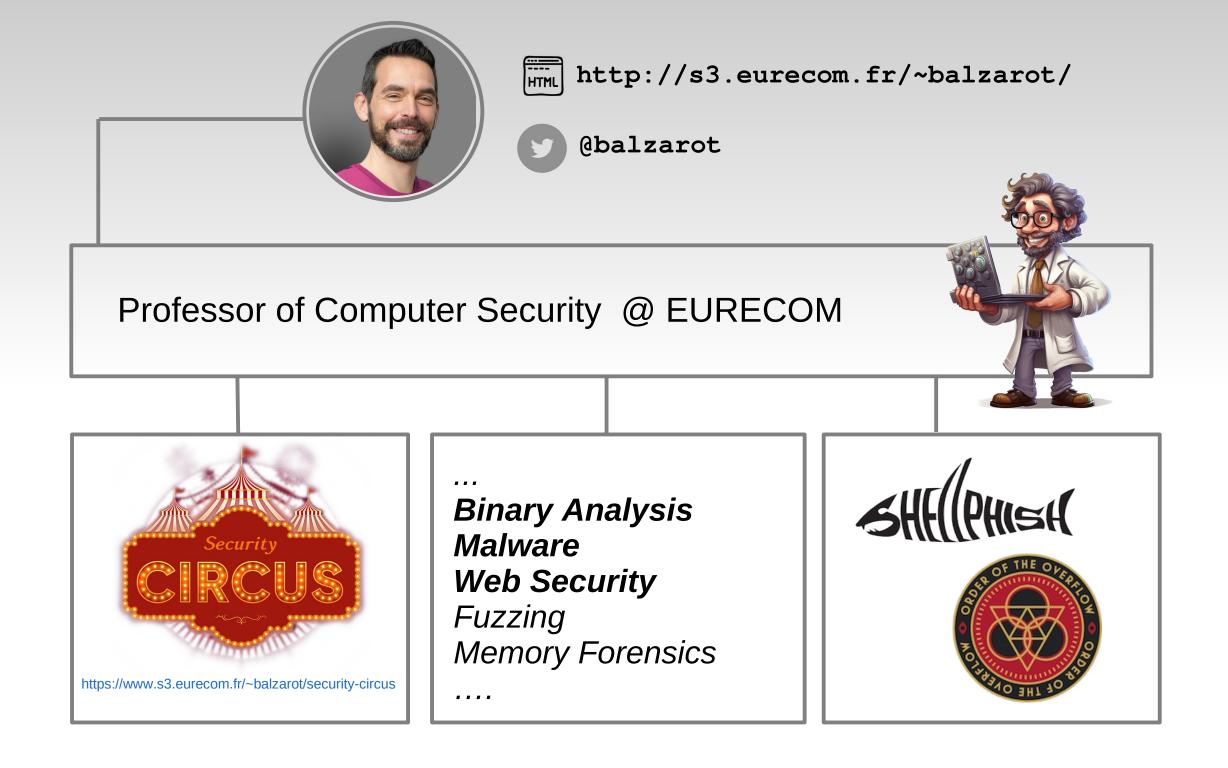
Davide Balzarotti



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Binary Analysis Malware Web Security Fuzzing Memory Forensics

. . . .





Mariano Graziano

Andrea Oliveri



$\texttt{fo} \cdot \texttt{ren} \cdot \texttt{sic}$

Adjective: Of, relating to, or denoting the application of *scientific methods and techniques* to the *investigation* of crime

Memory forensics

The preservation, collection, validation, identification, analysis, interpretation, documentation, and presentation of digital evidences extracted from the memory



(my definition) Memory forensics:

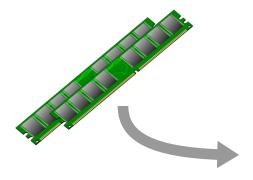
Reverse Engineering on Steroids

Pros

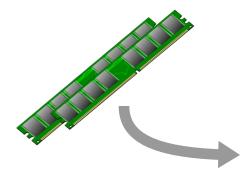
- Attackers often overlook their memory footprint
- Many of the kernel artifacts can be used for forensics
- Even rootkits designed to hide data in a running system need to be located somewhere in memory
- Certain information (loaded kernel modules, open sockets, ...) may be difficult to extract otherwise
- Some malware samples only reside in memory

Cons

- Memory is difficult to acquire
- The content of the memory keeps changing so even consecutive image acquisitions give different results
- Data collection requires an efficient approach with a small footprint
- Data structures change among different OSs and OS versions



751539f17641f7f15e5f5dc38d7426001f8945ec75513b4d89f28b75f029ce19c4205e5f5dc366900031d2f7f189c189eba58db600000000c38db42600000000ecc745f020000000



 75
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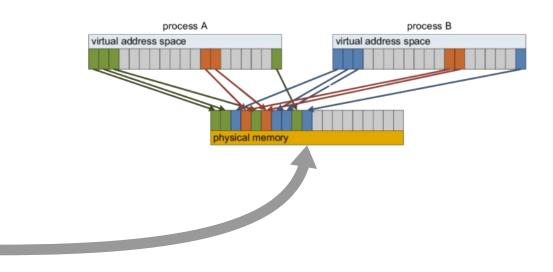
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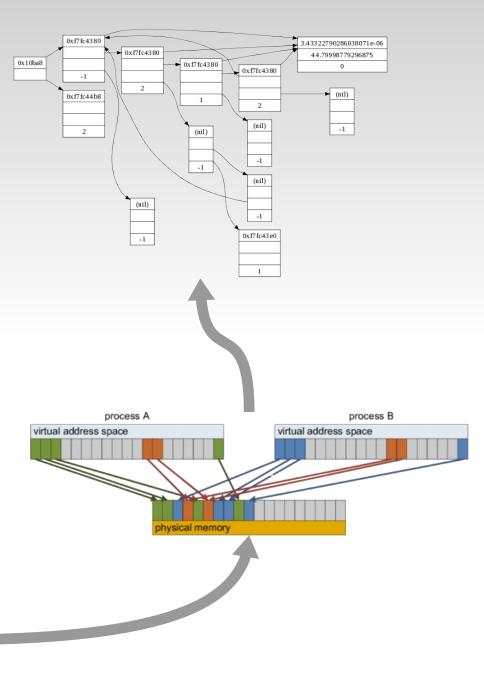
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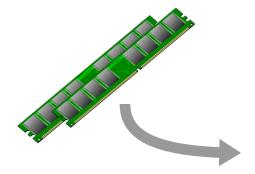
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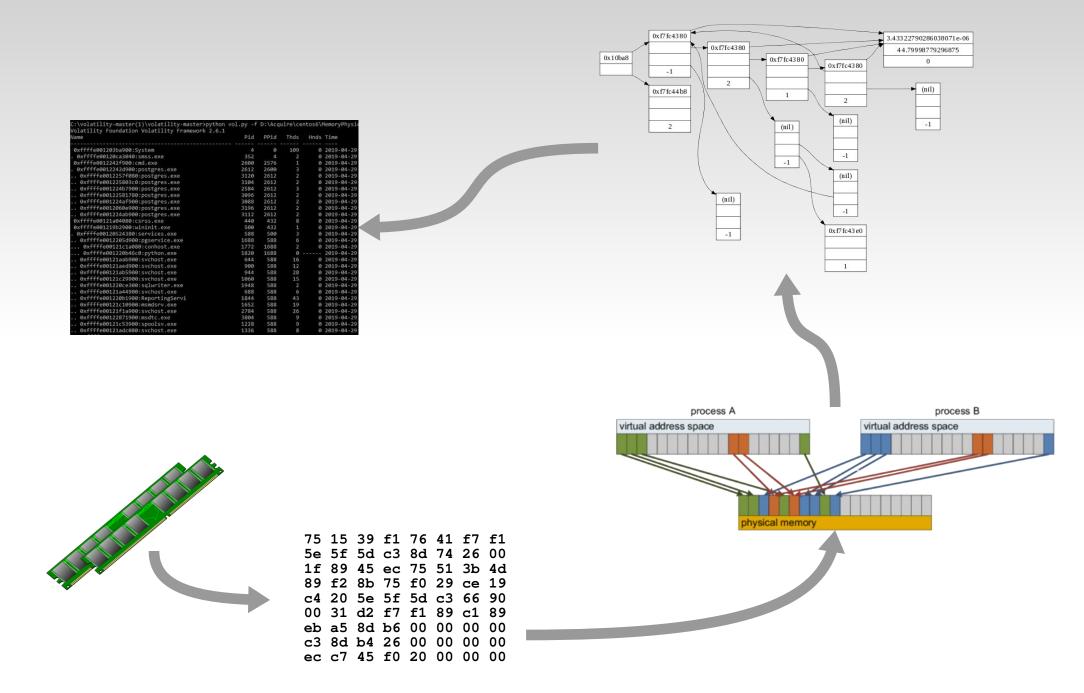
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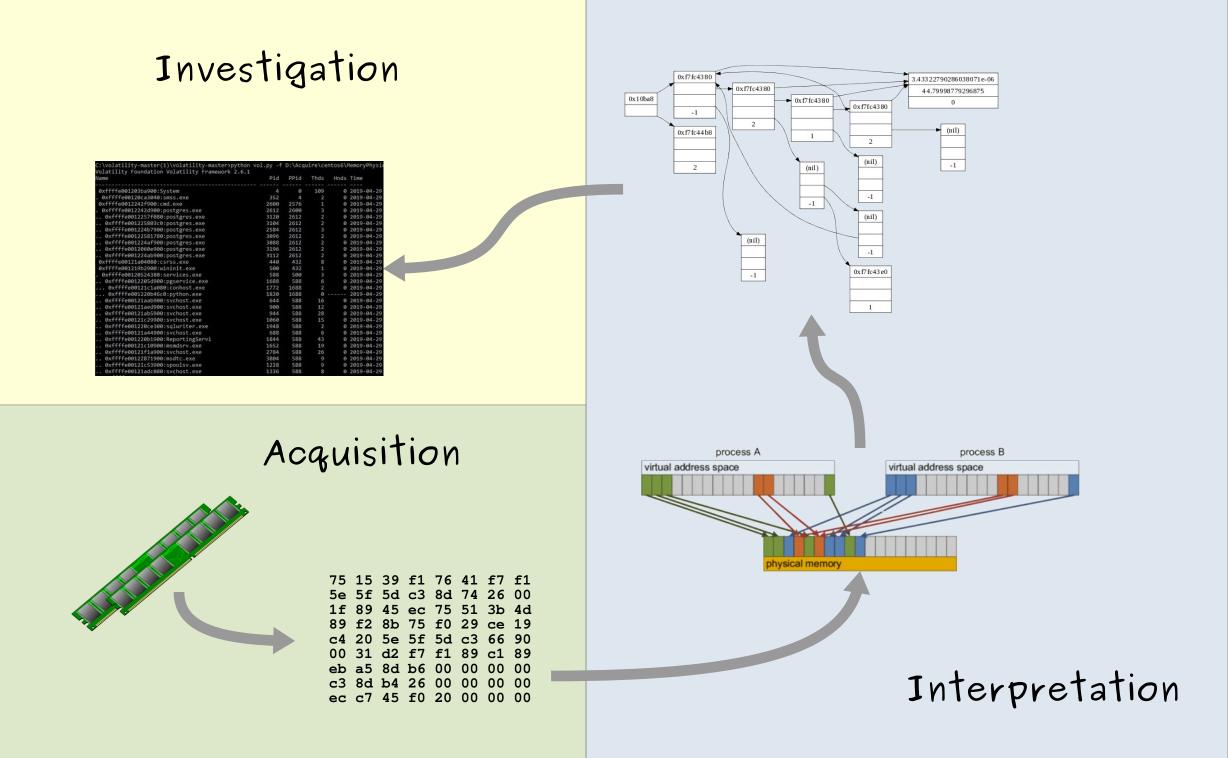
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Investigation

<u>چ</u>ھ

How to traverse data structures to recover high-level information

Acquisition

How to acquire a faithful copy of the physical memory

How to recover layout, location, and semantics of key data structures

Interpretation



(pre-2005) Carving

Memory Forensics 0.1

Looking for something you do not know in something you know

Looking for something you know in something you do not know

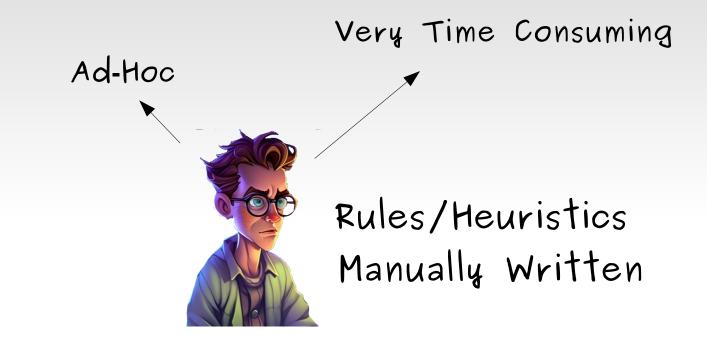
Looking for something you do not know in something you know – Structured Data Analysis –

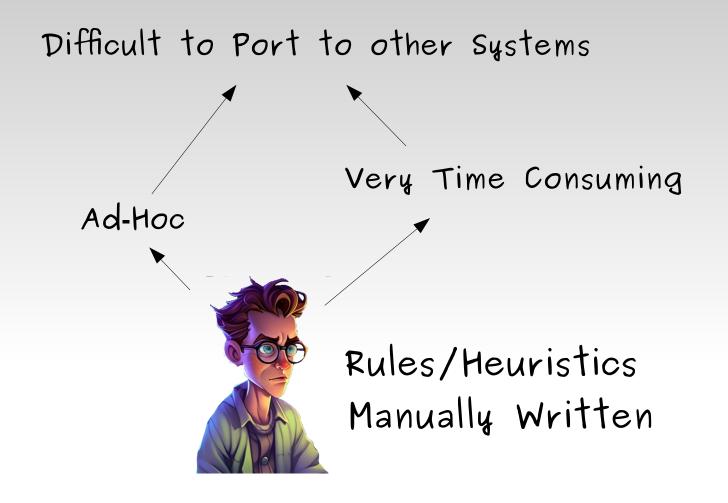
Looking for something you know in something you do not know – Carving –

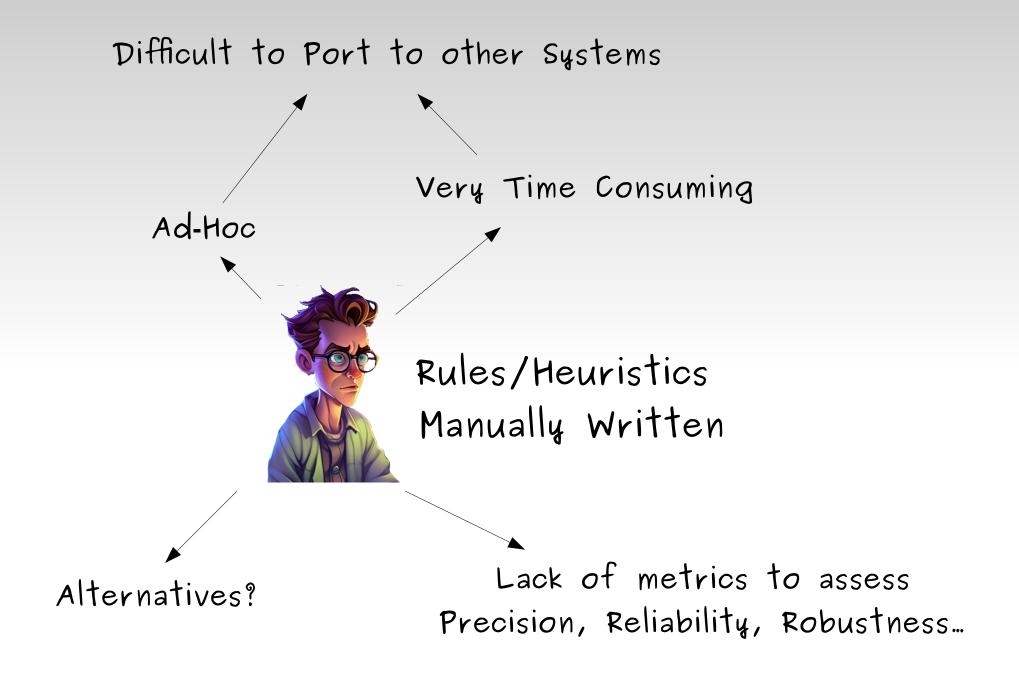


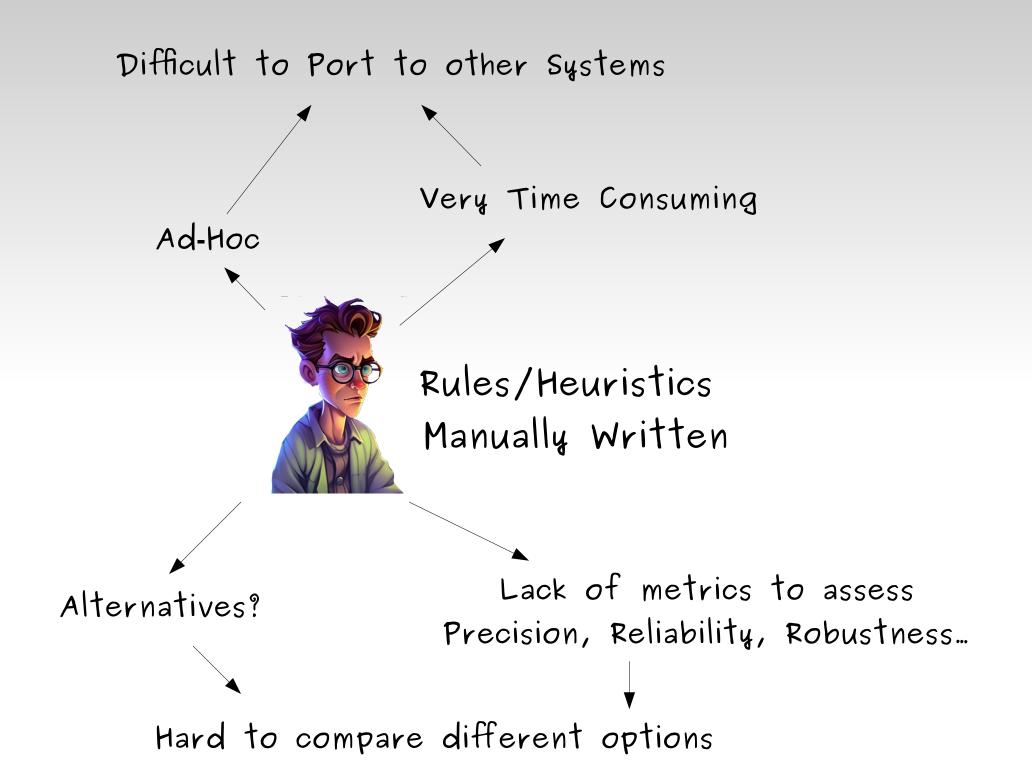
Rules/Heuristics Manually Written

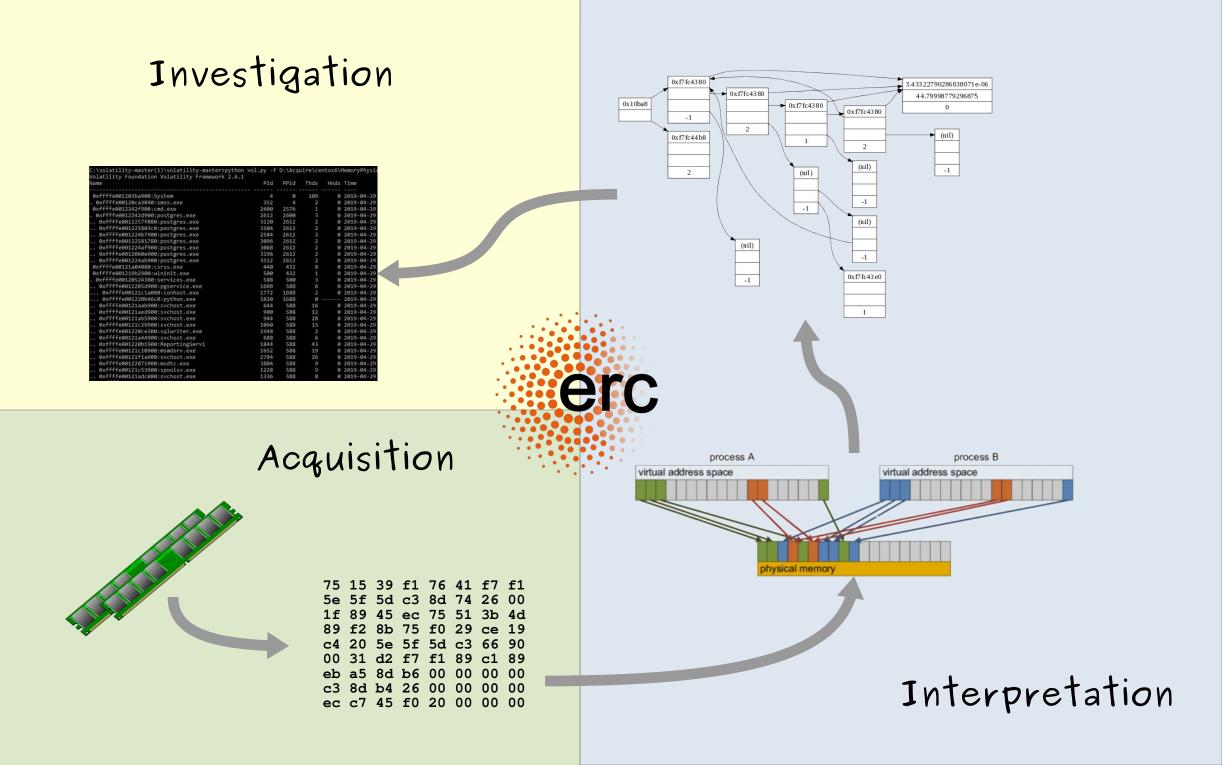
Memory Forensics 1.0

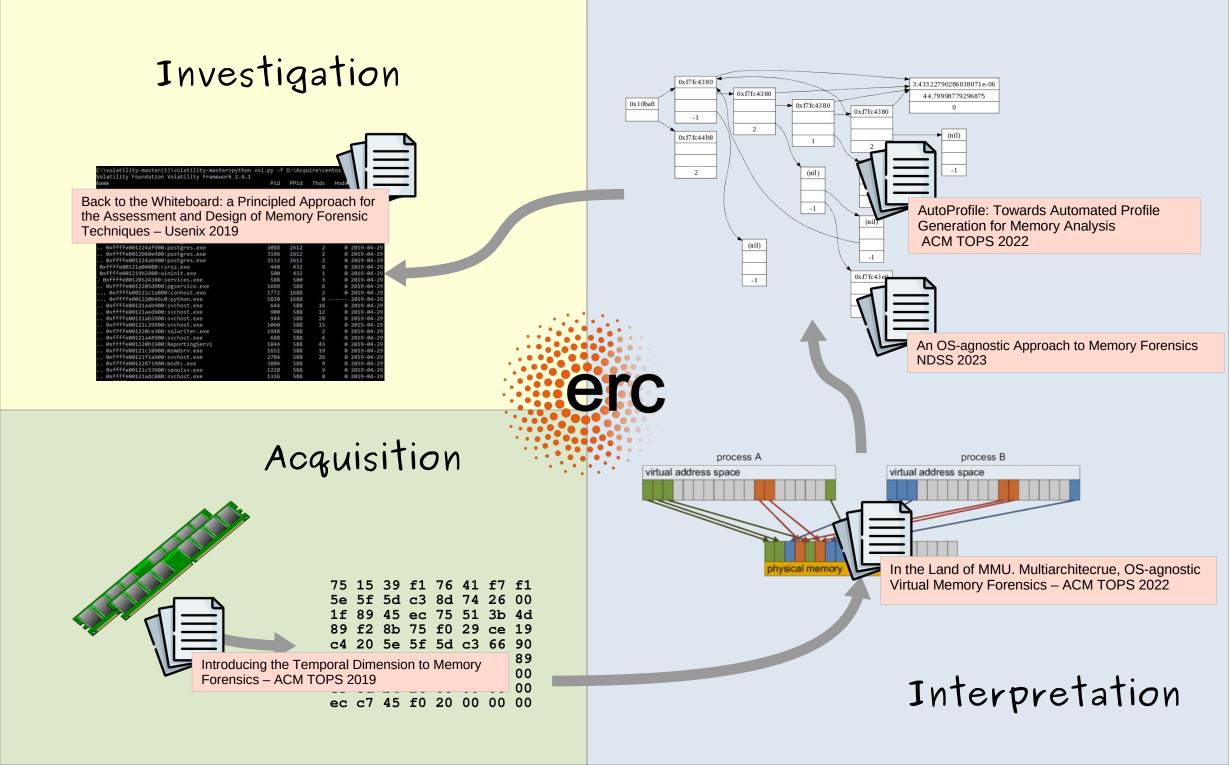


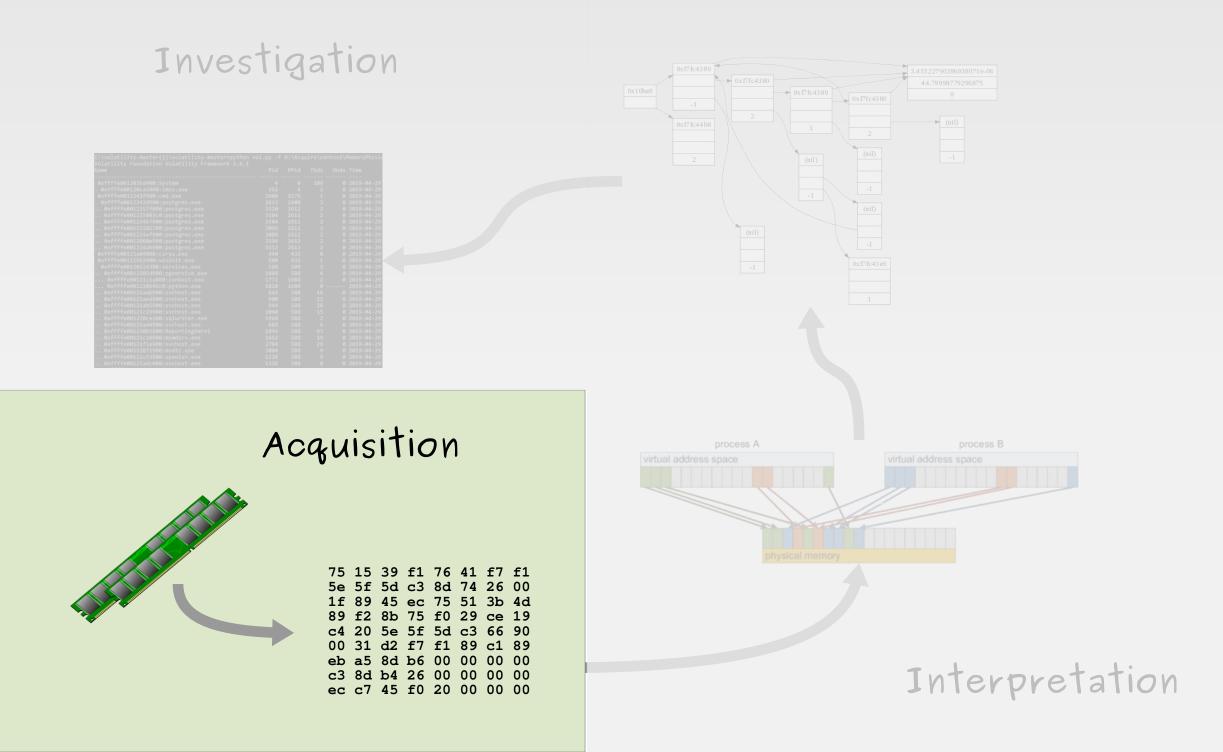












The physical address space is NOT contiguous: > sudo cat /proc/iomem

Hardware peripherals map registers or parts of their integrated memory into the physical address space via **Memory Mapped I/O**

Any attempt to read the memory mapped to a device would probably crash the system

OxFFFFFFFF..... 0xF080C000 PCI MMIO OxF03FFFFF. 0xF0020000. PCI MMIO 0xF0000000 0xE8000000 PCI MMIO 0xE0000000 ACPI Tables 0x7FFF0000. Memory 0x00100000 Upper BIOS 0x000F0000 Lower BIOS 0x000E0000 PCI Option ROMs 0x000000000 Video Window 0x000A000x0 EBDA 0x0009FC00 Memory 0x000000000

Software Acquisition

• Use software to read and dump the memory from within the system

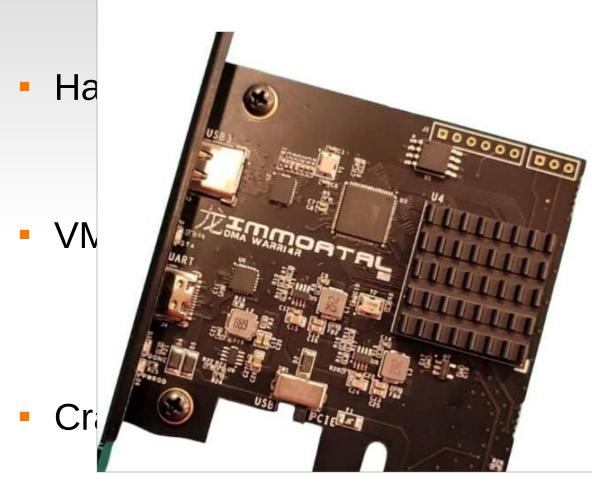
Hardware Acquisition

- Access memory from DMA
- Firewire, PCI-Express, USB 4, Intel DCI, Jtag

VM Acquisition

- Atomic acquisition
- New technologies like AMD Secure Encrypted Virtualization can block any type of memory dump from the hypervisor
- Crash dumps, hybernation files, ..
- Cold boot attacks

Software Acquisition



Cold boot attacks

Immortal DMA Warrior, FPGA DMA with Custom Unique PCILeech Firmware up to 275 MB/s Speed, FPGA DMA USB-C/PCIe Connection, FPGA USB Firmware Flash Capable, PCILeech DMA, Development Board, DMA, FPGA Brand: IMMORTAL DMA 5.0 ***** 1 rating

Currently unavailable.

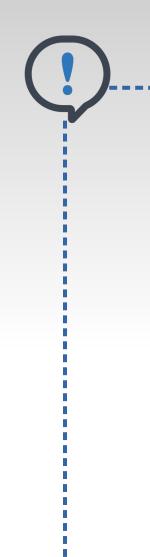
We don't know when or if this item will be back in stock.

Brand	IMMORTAL DMA							
Hardware	USB, PCI							
Interface								
Style	Classic							

About this item

đ

 Pre-Flashed Individual Custom Firmware (PCILeech): Firmware customized to prevent detection from some of the toughest anticheats and malware. Each individual customized firmware of PCILeech is destroyed aftering being flashed to your FPGA DMA device to guarantee individuality.



A complete memory acquisition takes several minutes, during which the OS is running

The Problem of (lack of) Atomicity

A complete memory acquisition takes several minutes, during which the OS is running

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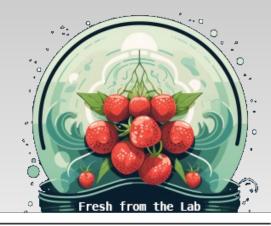
When idle, the Linux kernel performs over 300K write operations per second

The Problem of (lack of) Atomicity

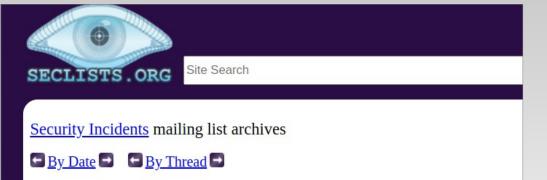
A complete memory acquisition takes several minutes, during which the OS is running

When idle, the Linux kernel performs over 300K write operations per second





Mode	Writes on kernel address space (Millions)		Writes on MMIO regions		Total size (GiB)		Unique physical pages		Time required (ratio)	
	USB	SATA	USB	SATA	USB	SATA	USB	SATA	USB	SATA
Btrfs	874	811	59824	37778	6.01	5.58	249340	249204	1.81x	1.53x
exFAT	1005	938	96112	61772	6.89	6.43	251074	250728	1.79x	1.33x
Ext4	818	757	60692	35696	5.61	5.20	249421	248873	1.76x	1.16x
Ext4 no journal	776	719	61744	36443	5.33	4.95	249439	248864	1.78x	1.10x
F2FS	951	910	61329	36743	6.51	6.23	249406	249379	2.04x	1.33x
NTFS	796	739	61329	38711	5.48	5.09	249411	249129	1.75x	1.31x
FAT32	1404	1317	84456	89542	9.65	9.06	250328	250908	2.39x	1.82x
XFS	632	569	57605	34255	4.41	3.97	249405	249041	1.62x	1x
Btrfs D. I/O	49137	37708	9147061	6707022	344.04	265.19	75037	78885	255.76x	73.00x
exFAT D. I/O	10698	4713	5204034	3950081	73.20	32.11	466	497	109.34x	15.85x
FAT32 D. I/O	16657	6000	9138277	5773820	114.15	40.88	1125	1127	236.71x	19.58x
Network	1336	-	1000453	-	8.64	-	488	-	2.73x	-



List Archive Search

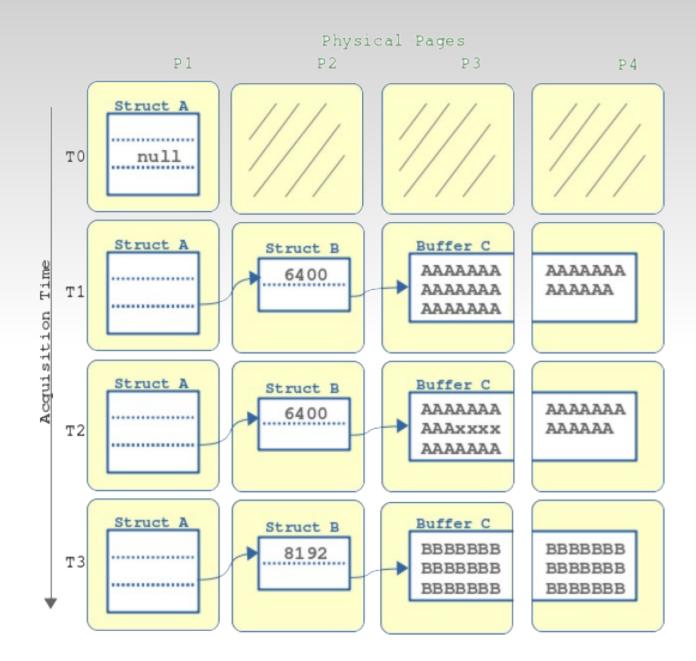
Re: Digital forensics of the physical memory

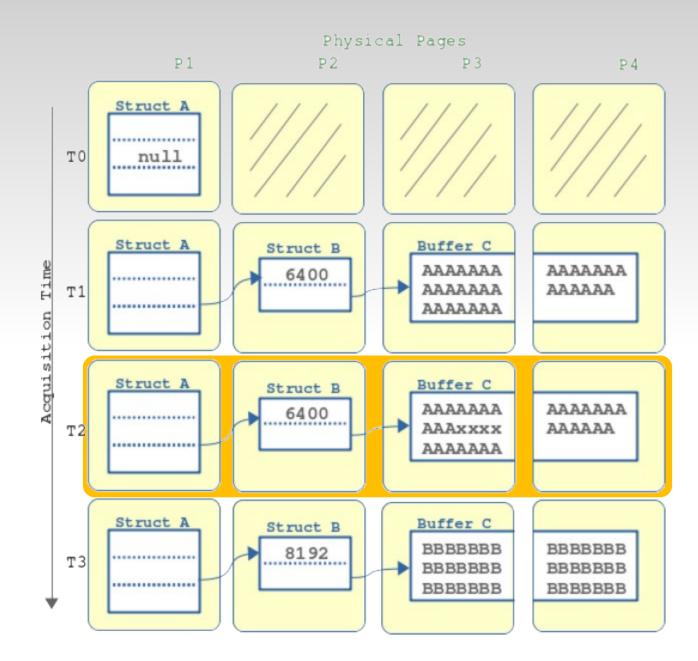
From: Harlan Carvey <keydet89 () yahoo com> Date: Fri, 17 Jun 2005 09:35:16 -0700 (PDT)

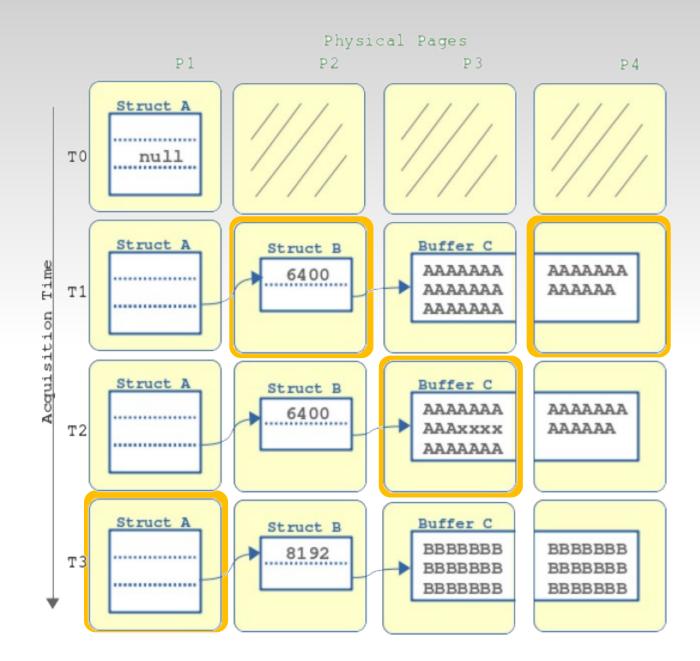
One of the issues in particular is that he starts off by mentioning the FU rootkit and the SQL Slammer worm, both of which are specific to Windows...and then presents examples using only a Linux system. He states in the paper that similar work can be done on Windows systems, but never provided any information to that effect.

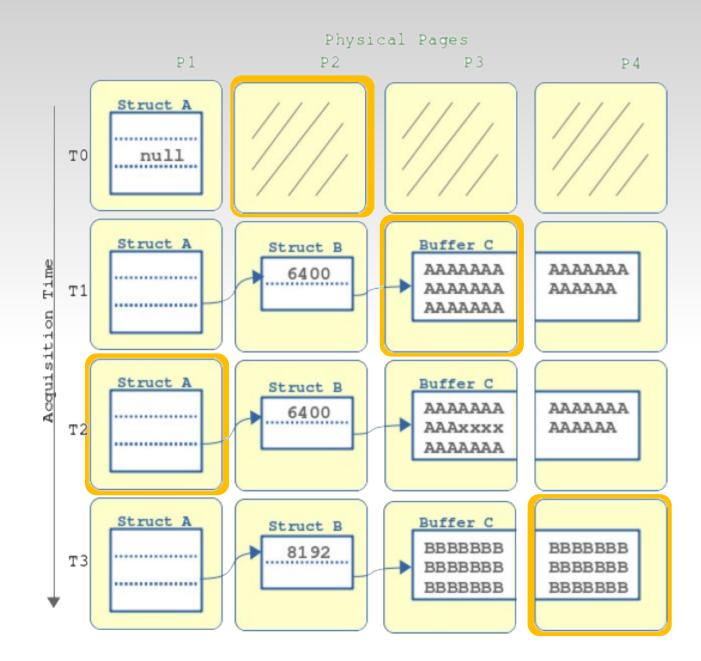
Based on entries I made to my blog the other day, I ended up having a conversation w/ someone from MS about this very issue. The issue of using dd.exe to image Physical Memory goes beyond the fact that there don't seem to be any maps describing how physical memory is used by Windows systems, and that memory used by processes consists of both RAM and the pagefile. Additional issues include, as you pointed out, that while the imaging process is occurring, the kernel memory (and even user-mode memory) is changing...so what you end up with is a smear, for want of a better term.

Even tools like pmdump.exe and LiveKD (SysInternals.com) are not sufficient for collecting user-mode memory, b/c they do not lock or suspend memory.











Introducing the Temporal Dimension to Memory Forensics

\$./vol.py -f dump.raw --profile=... --pagetime pslist
<original pslist output>

```
Accessed physical pages: 171
Acquisition time window: 72s
[XX-----XxX----XXXX---XX---Xxx-X-X-XXX]
```



Ongoing experiments repeated on 10 dumps

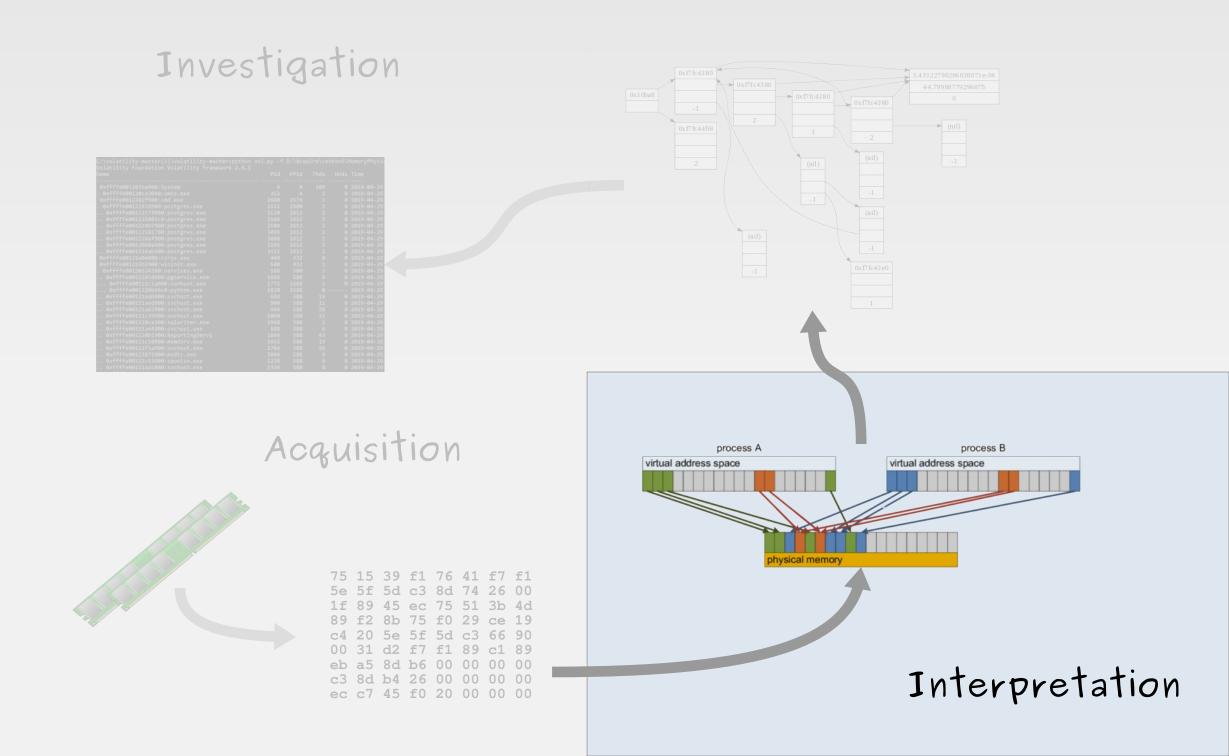


ALL contain inconsistencies in page tables

The kernel is ALWAYS affected

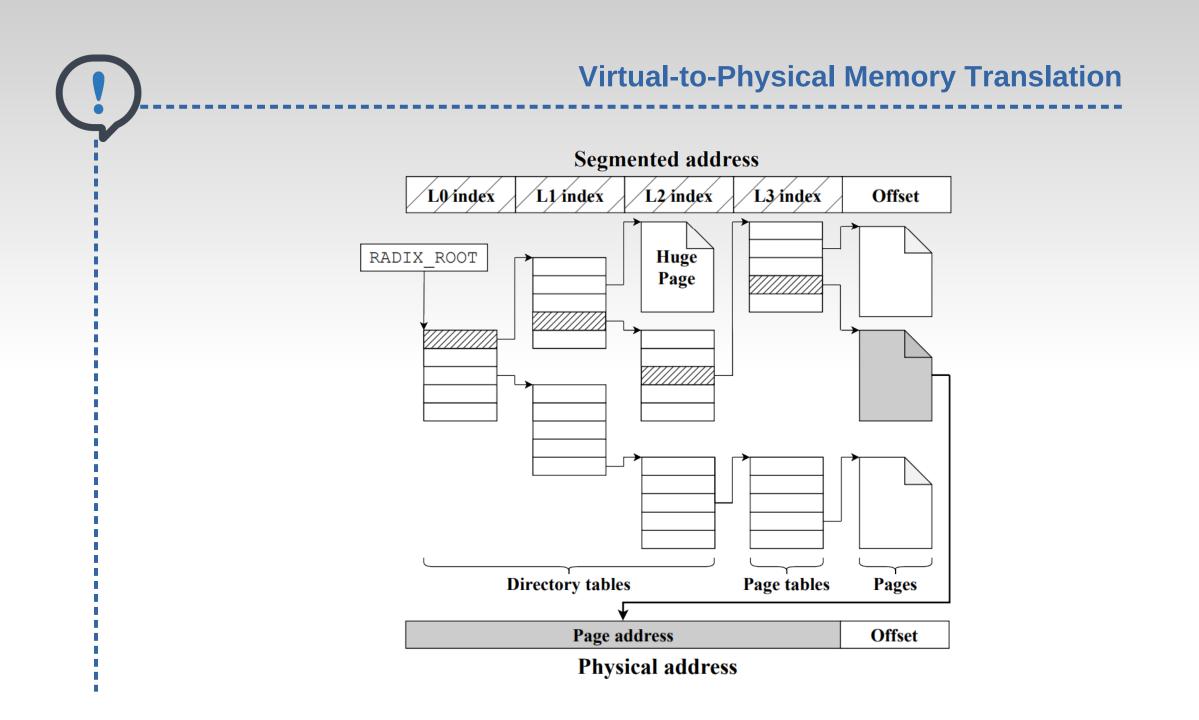
Dozens of processes with corrupted address spaces

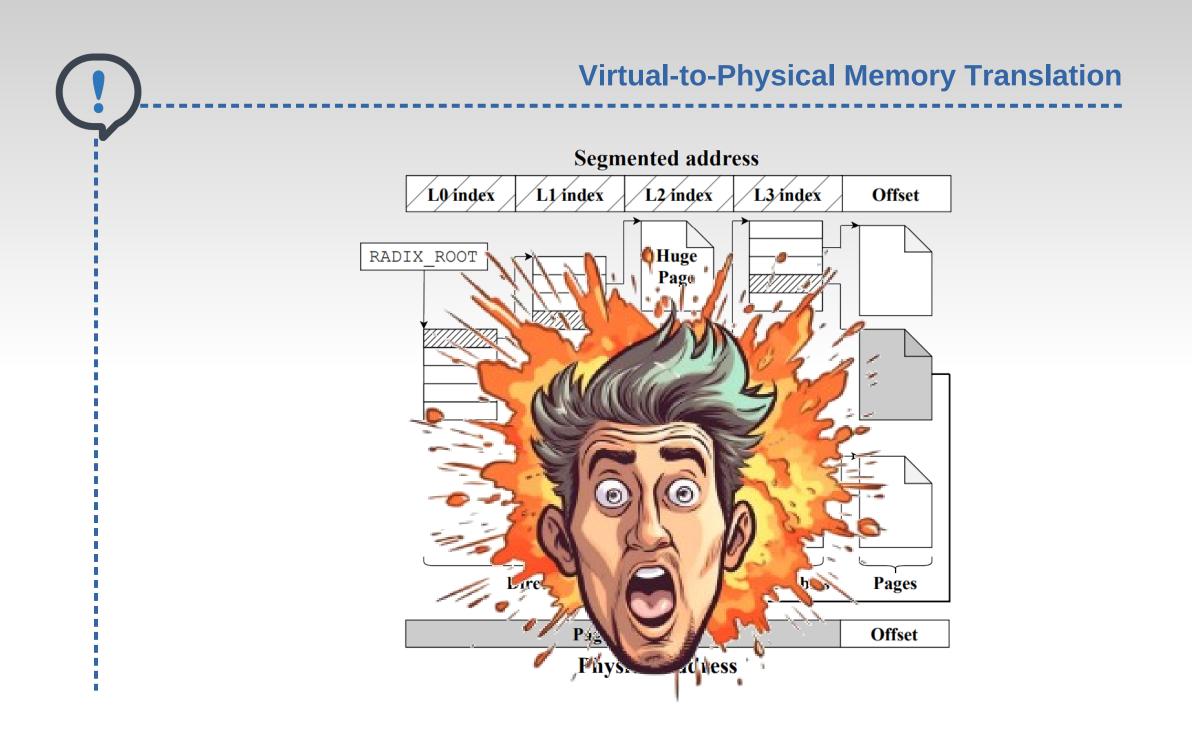
Two cases in which the pages of one process get attributed to another

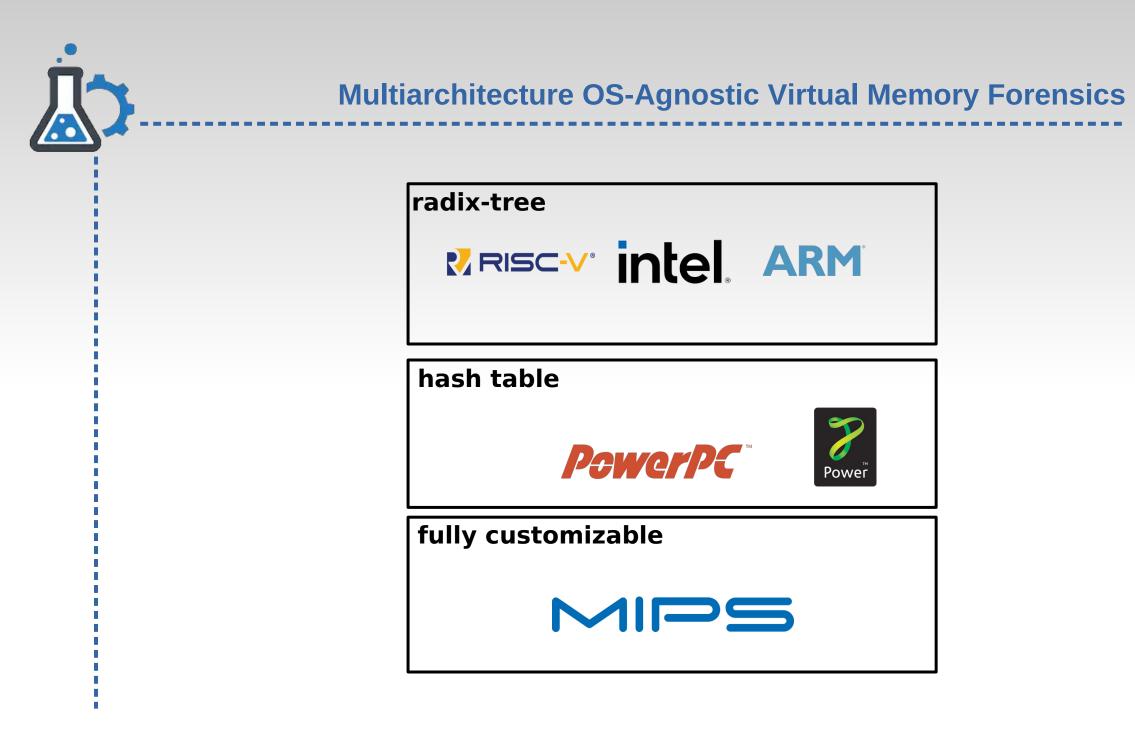


The V2P translation is performed in hardware by the Memory Management Unit (**MMU**) based on in-memory data structures and dedicated CPU registers that are configured by the OS

The translation process can involve **segmentation** and **paging**. Some architectures use one or the other, some use both.







Multiarchitecture OS-Agnostic Virtual Memory Forensics

Want to know more?

In the Land of MMUs: Multiarchitecture OS-Agnostic Virtual Memory Forensics

ANDREA OLIVERI, Eurecom, France DAVIDE BALZAROTTI, Eurecom, France

The first step required to perform any analysis of a physical memory image is the reconstruction of the virtual address spaces, which allows translating virtual addresses to their corresponding physical offsets. However, this phase is often overlooked and the challenges related to it are rarely discussed in the literature. Practical tools solve the problem by using a set of custom heuristics tailored on a very small number of well-known operating systems running on few architectures.

In this paper, we look for the first time at all the different ways the virtual to physical translation can be operated in 10 different CPU architectures. In each case, we study the inviolable constraints imposed by the MMU that can be used to build signatures to recover the required data structures from memory without any knowledge about the running operating system. We build a proof-of-concept tool to experiment with the extraction of virtual address spaces showing the challenges of performing an OS-agnostic virtual to physical address translation in real-world scenarios. We conduct experiments on a large set of 26 different OSs and a use case on a real hardware device. Finally, we show a possible usage of our technique to retrieve information about user space processes running on an unknown OS without any knowledge of its internals.

Additional Key Words and Phrases: memory forensics, OS-agnostic forensics, virtual memory, MMU

ACM Reference Format:

Andrea Oliveri and Davide Balzarotti. 2022. In the Land of MMUs: Multiarchitecture OS-Agnostic Virtual Memory Forensics. 1, 1 (April 2022), 33 pages. https://doi.org/10.1145/nnnnnnnnnnn

1 INTRODUCTION

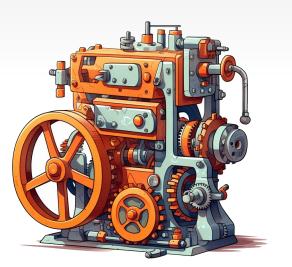
The problem of recovering semantic information from low-level data is common to many areas of computer security. In particular, this is the main obstacle when performing a physical memory analysis—a task that is key for both memory forensics and virtual machine introspection. The problem, often called the *semantic gap*, captures the challenge of "interpreting low level bits and bytes into a high level semantic state of an in-guest operating system" [35]. However, at a closer look, the semantic gap can be further divided into two different aspects: the reconstruction of the virtual address spaces (which deal with translating pointers expressed as virtual addresses to their physical position in the



- 1. Structural Signatures derived by inviolable MMU constraints
- Validation Rules based on inviolable constraints imposed by other CPU subsystems (e.g., pages containing the Interrupt Address Table should be mapped in all VASs)
- 3. Binary code analysis to recover MMU-related CPU registers



Physical to Virtual Memory Translation



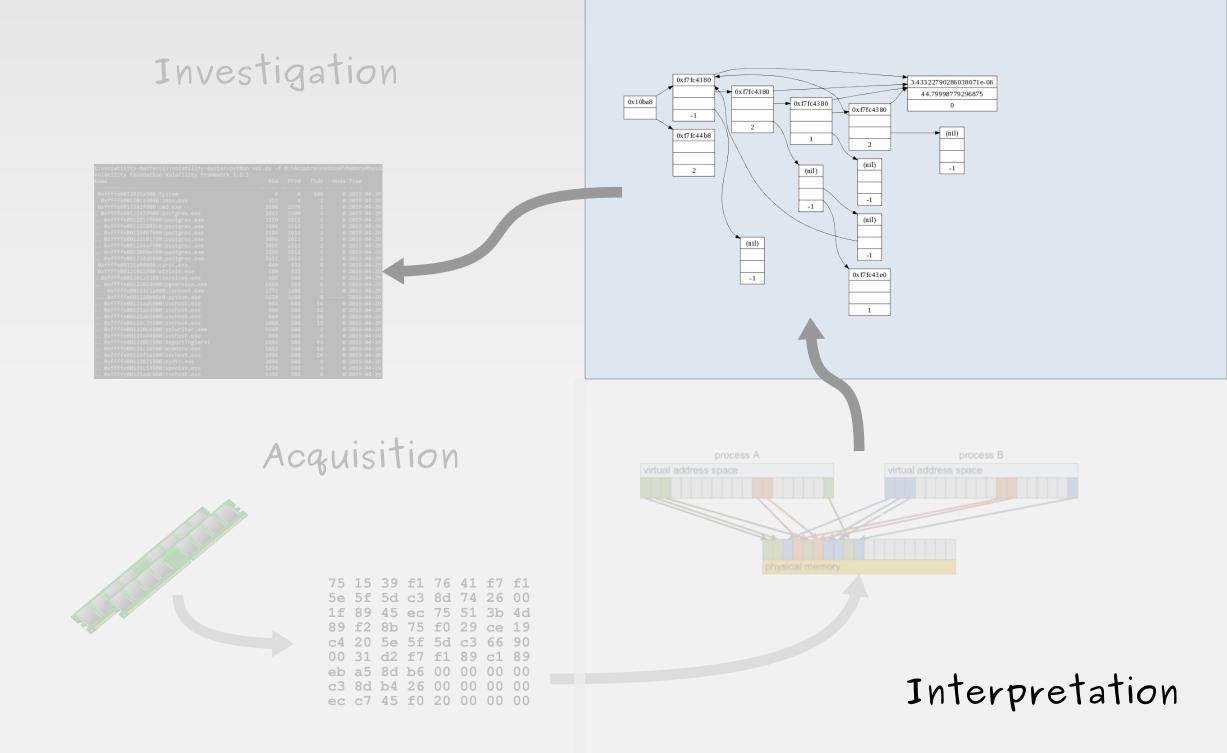
MMUShell

https://github.com/eurecom-s3/mmushell



OS		Architectures MMU modes				
	Open-source Kernel type ¹	x86 PAE x86 IA32 RISC-V SV48 RISC-V SV32 PowerPC MIPS32 Radix MIPS32 TLBs ARM32 Short ARM32 Short AMD64				
9Front[24]	Н●	•				
Barrelfish[17]	U •					
Darwin[4]	H •	•				
Embox[5]	R •	• • • •				
FreeBSD	M •	•• ••				
GenodeOS[6]	m 🔴					
HaikuOS[7]	Н●	• •				
HelenOS[8]	m 🔴	••••				
Linux Buildroot[3]	M •					
Linux Debian	М •					
MacOS 9	n 🔾	•				
MacOS X	H O	•				
Minix3[9]	m 🔴	•				
MorphOS[10]	m 🔾	•				
NetBSD	M •	••• • •				
Illumos[29]	M •	•				
QNX[11]	RO	•				
rCore[13]	M •	•• ••				
ReactOS[14]	m 🔴	•				
RedoxOS[15]	m 🔴	•				
vxWorks[19]	R 🔾	•				
Windows 10	H O	••				
Windows 95	M ()	•				
Windows NT	H O	•				
Windows XP	H 🔾	• ••				
XV6[20]	М •	•				

ory Translation



Memory Forensics is based on **PROFILES**, which contain precise descriptions of all the kernel data structures necessary to perform the analysis.

Q1: Can we automatically generate profiles starting from the dump itself?

Q2: Can we perform some analysis also *without any* profile?





The Problem with Profiles

The important is NOT how much kernel structures change across kernels

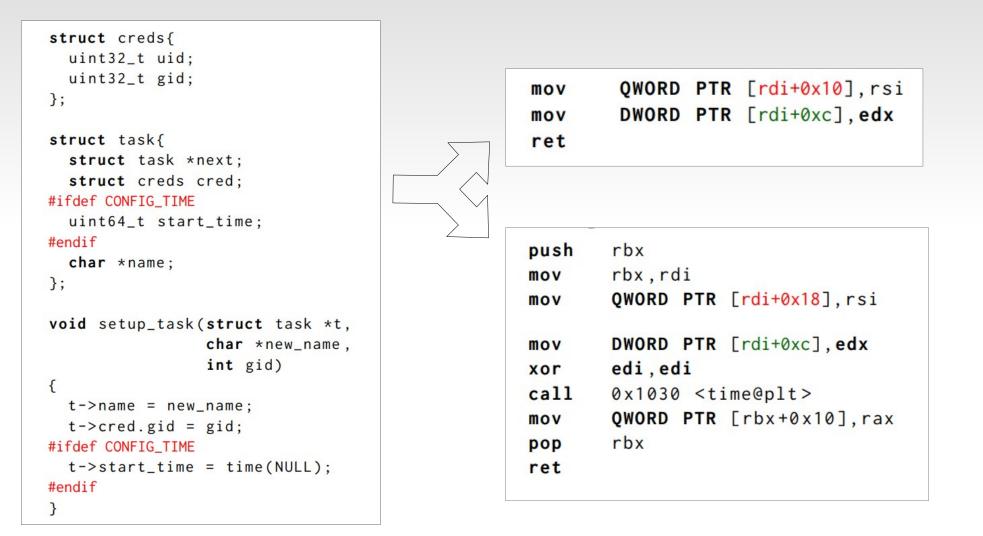
But how much they change within a single version – because of user configurations or compiler options. E.g., The layout of task_struct is shaped by more than 60 different #ifdef

Modern kernels also support *structure layout randomization* as a form of protection against exploitation

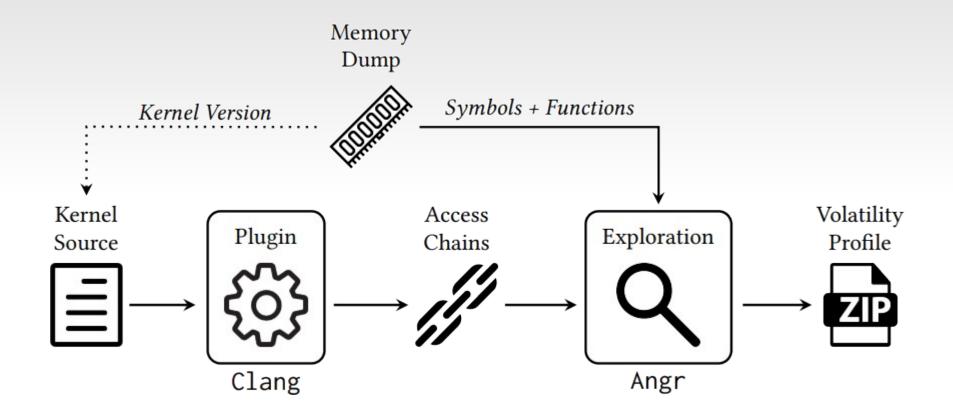


While the struct definitions are lost during the compilation process, they are "reflected" in the code itself.









*AutoProfile: Towards Automated Profile Generation for Memory Analysis



Version	Release Date	Configuration	Used Fields	Extracted Fields	
4.19.37	04/2019	Debian	234	220 (94%)	
4.19.37	04/2019	Debian + RANDSTRUCT	234	194 (83%)	
5.6.19	03/2020	Raspberry Pi	227	217 (95%)	
4.4.71	06/2017	OpenWrt	236	216 (92%)	
3.18.94	05/2018	Goldfish (Android)	239	220 (92%)	
2.6.38	03/2011	Ubuntu	226	213 (94%)	



Katana (very very similar solution published one year later)

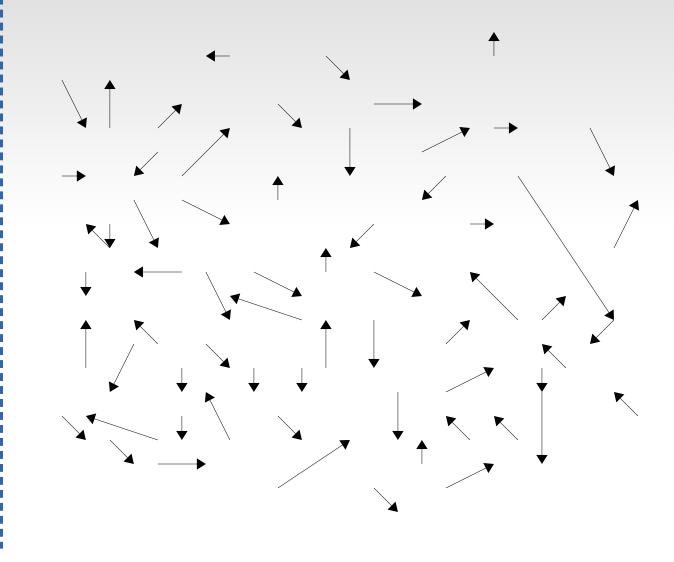
https://github.com/tum-itsec/katana

*Katana: Robust, Automated, Binary-Only Forensic Analysis of Linux Memory Snapshots

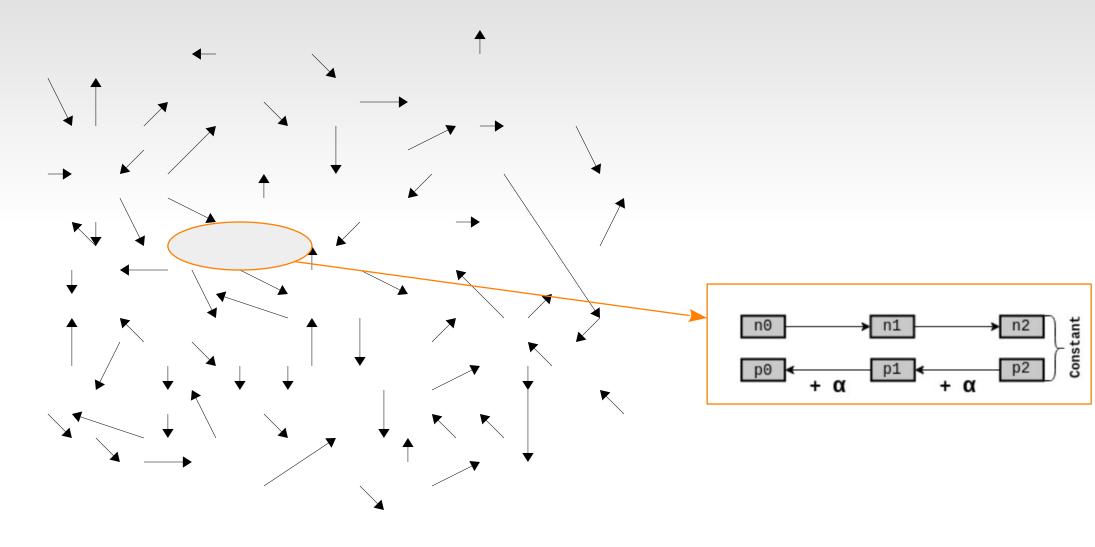


Look Mum, no Profiles!

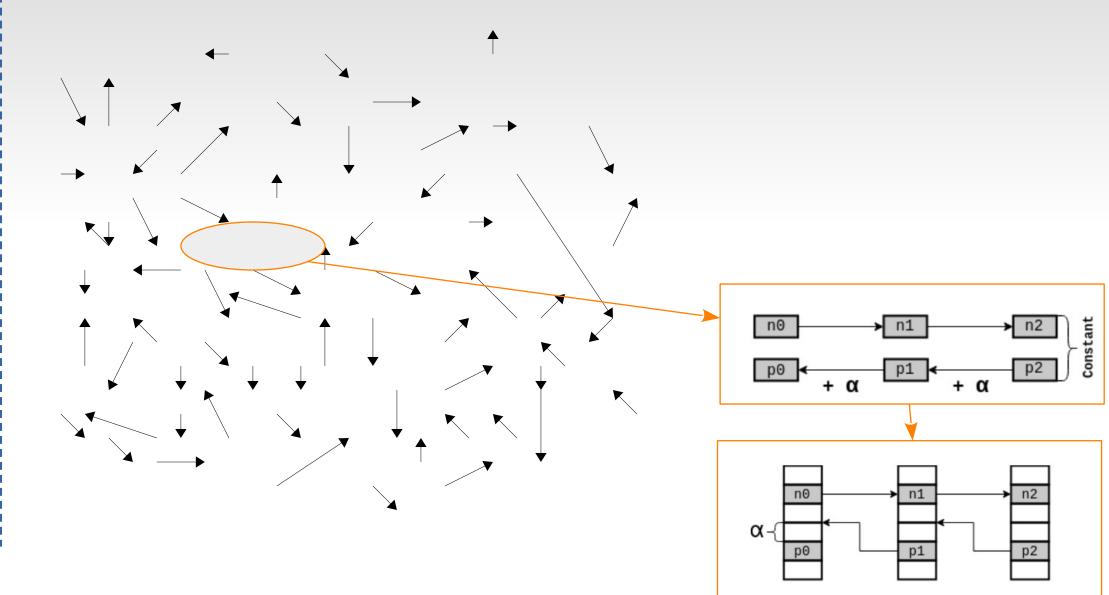




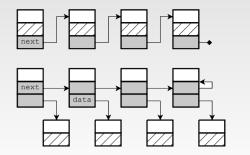


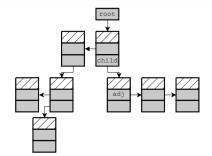


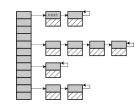


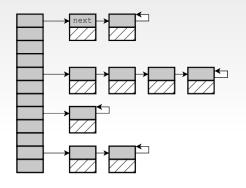


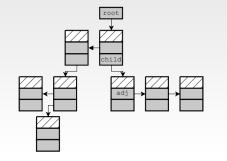


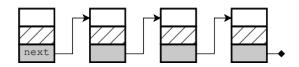


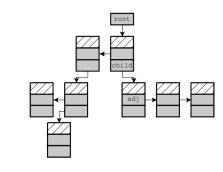












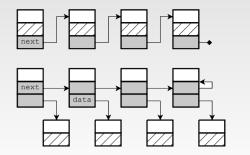


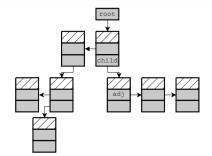
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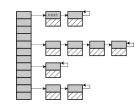
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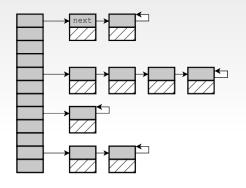
	OS	Linear DL. L.	Circular DL. L.	Trees	Arrays of *structs	Arrays of *strings	Linked Lists
	Darwin	11	385	127	1214	1801	35
root	Embox	0	22	35	1131	795	6
	FreeBSD	86	0	993	1008	895	41
child	HaikuOS	4117	64	0	305	232	1184
	HelenOS	25	1173	127	41	45	1
adj 🔶	iOS	20	256	192	5234	229	36
	Linux	120	3632	1034	693	5947	46
	Linux (Aarch64)	110	3362	936	229	4985	43
	NetBSD	41	18	1218	1482	406	45
	ReactOS	7	200	49	492	325	12
	ToaruOS	101	0	14	62	229	15
	vxWorks	51	14	199	349	416	13
	Windows XP	38	889	228	463	206	20
	Windows 10	145	6639	36	0	282	0

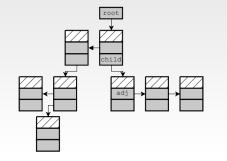


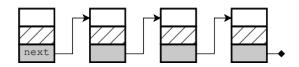


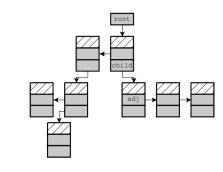




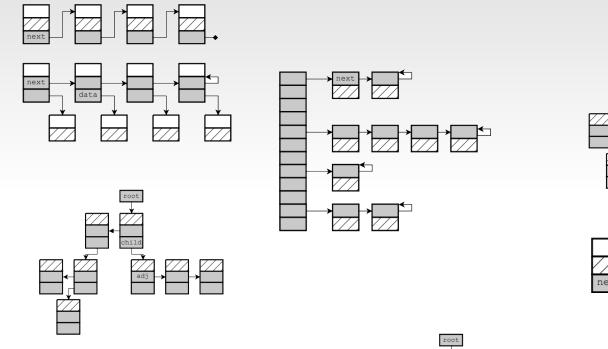


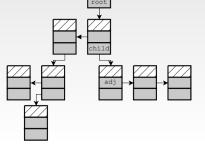


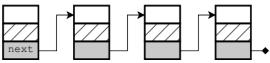


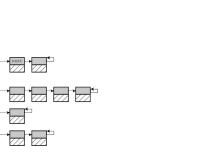


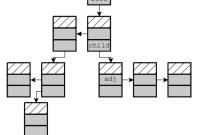






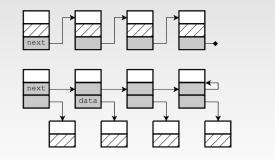


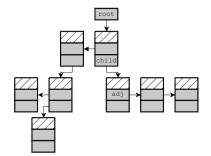


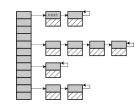


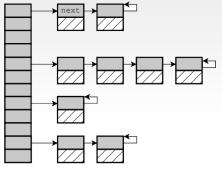


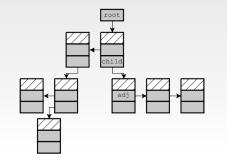


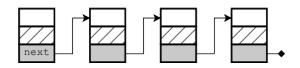


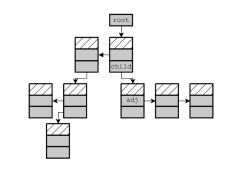






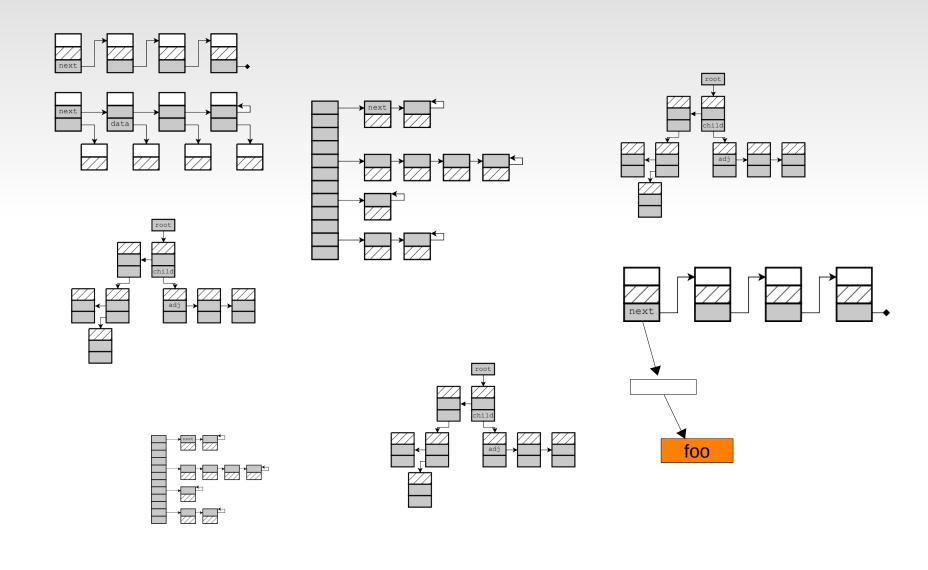




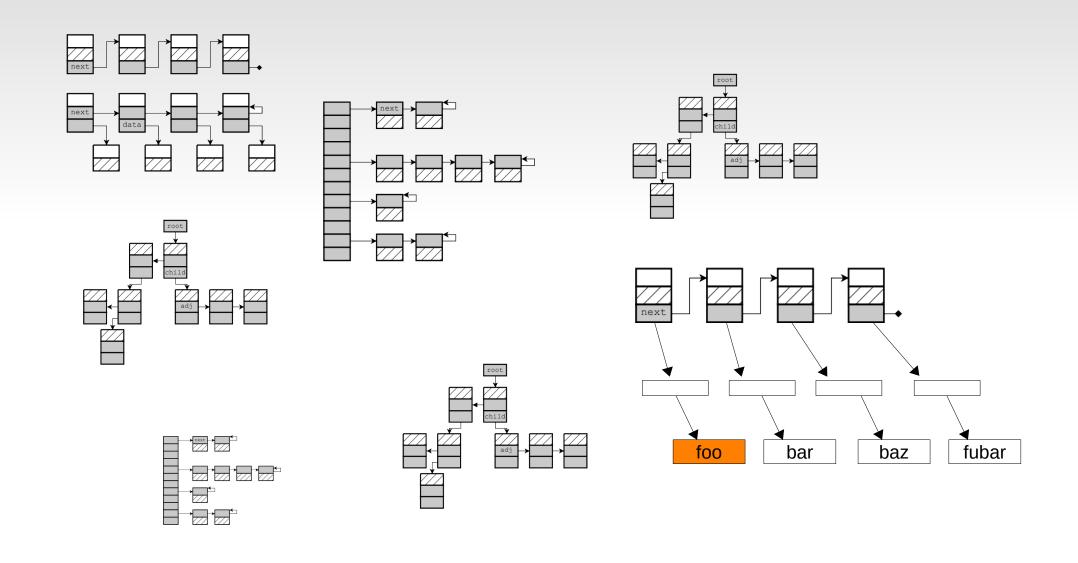




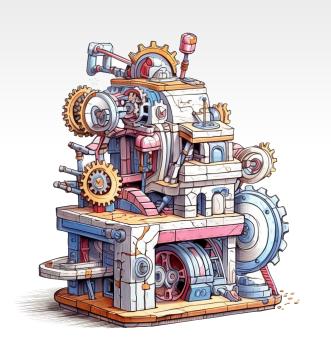












Fossil

https://github.com/eurecom-s3/fossil



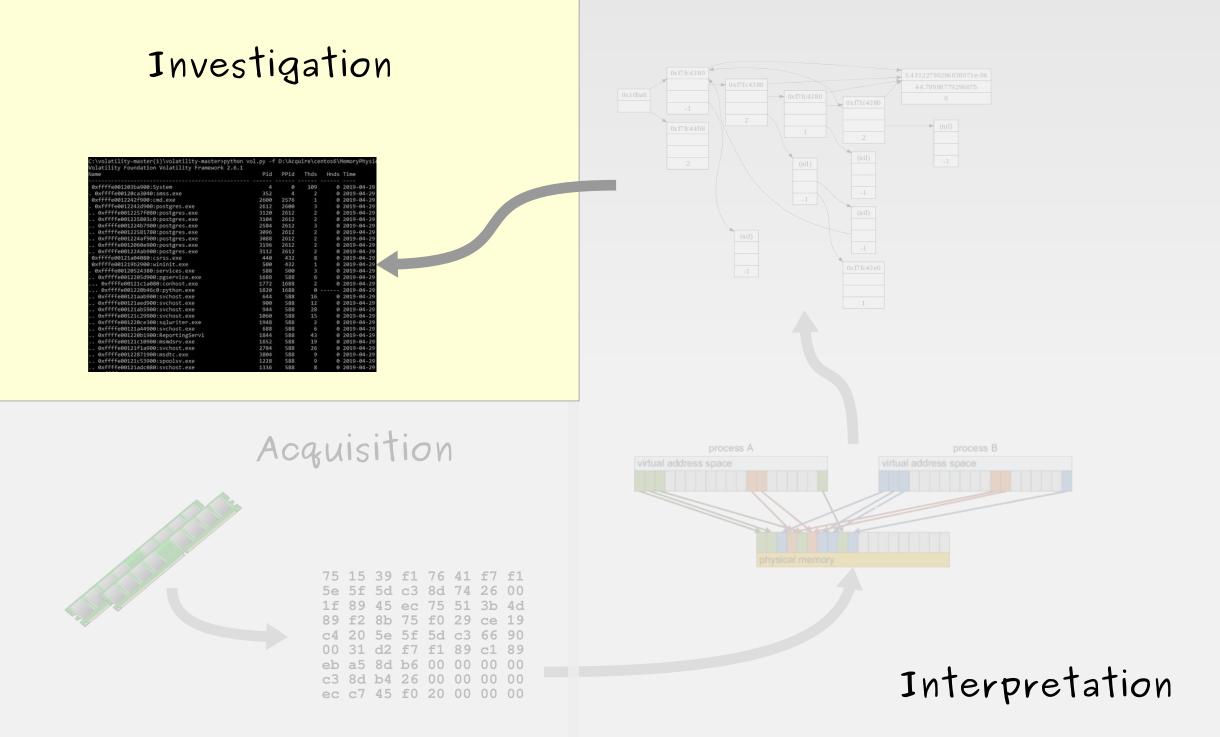
OS	Kernel modules	Kernel pools	File systems	Other structures
Darwin	•	•	•	• List of network devices • System locks • Ker- nel/user pipes • Kernel parameters
Embox	٠		0	List of commands
FreeBSD	٠		•	
HaikuOS	•	•	0	 Executable libraries Kernel/user pipes Semaphores
HelenOS	•	•	•	
iOS	0	•	•	• List of network devices • System locks • Ker- nel/user pipes • Kernel parameters
Linux	•	•	٠	• Files in sysfs • Network protocols
Linux (AArch64)	•	•	•	• Files in sysfs • Network protocols
NetBSD	•	•	٠	• Kernel tasks
ReactOS	0	•	0	
ToaruOS	•		•	Devices' list Processes' environment
vxWorks	0	•	0	Devices' list Open sockets
Windows XP	•	•	0	
Windows 10	•	•	•	





OS	Process list	Kernel modules	Kernel pools	Filesystem
Darwin	2	10	11	7
Embox	17	0		
FreeBSD	24	31		26
HaikuOS	6	1	11	
HelenOS	4	2	1	1
iOS	2		2	15
Linux	5	28	26	15
Linux (AArch64)	4	22	19	24
NetBSD	2	6	18	0
ReactOS	5		12	
ToaruOS	3	2	3	
vxWorks	4		2	
Windows XP	5	1	2	
Windows 10	41	0	0	0

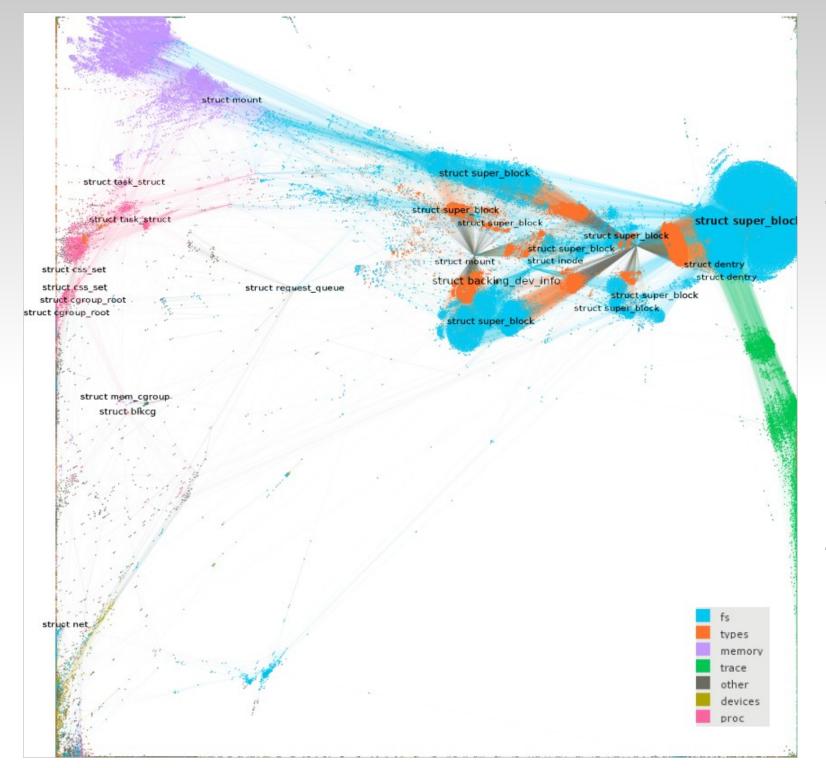






The goal of the analyst is to traverse the graph of kernel data structures to locate the information she needs.

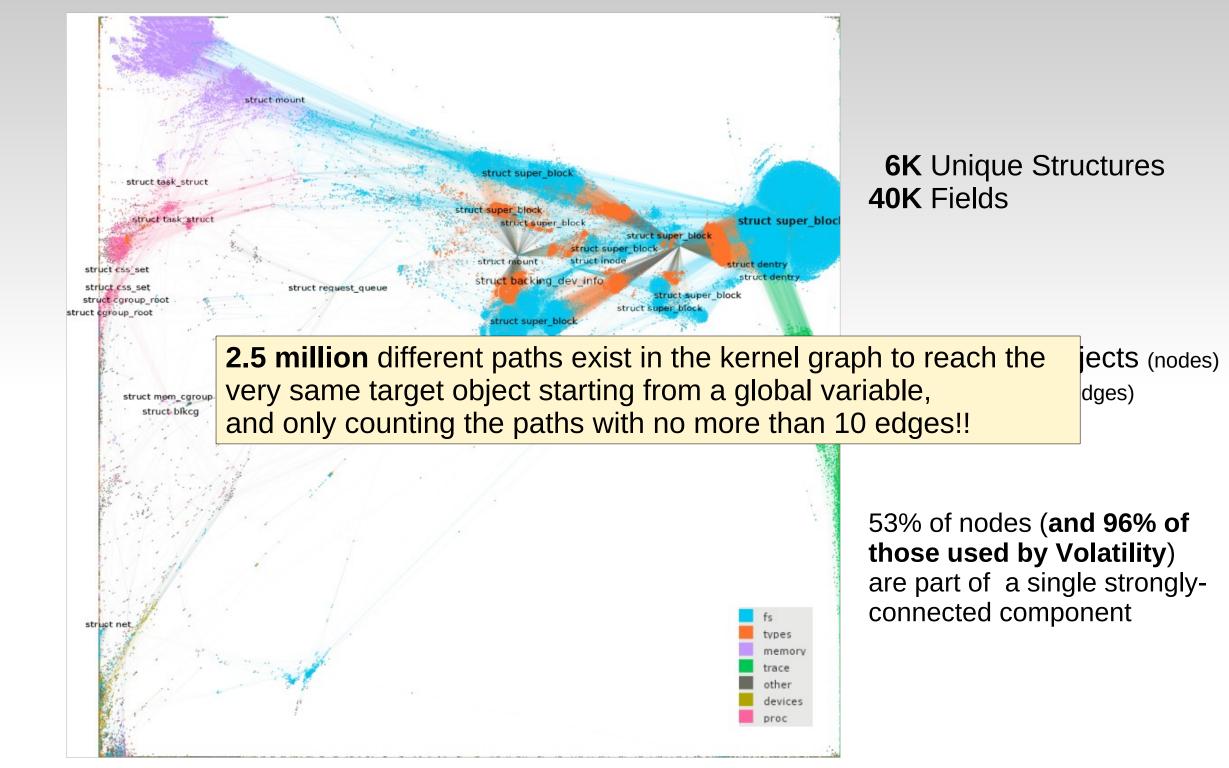
Each rule (e.g., a plugin to list processes) corresponds to a set of paths on the graph.



6K Unique Structures **40K** Fields

100K Kernel Objects (nodes) **840K** Pointers (edges)

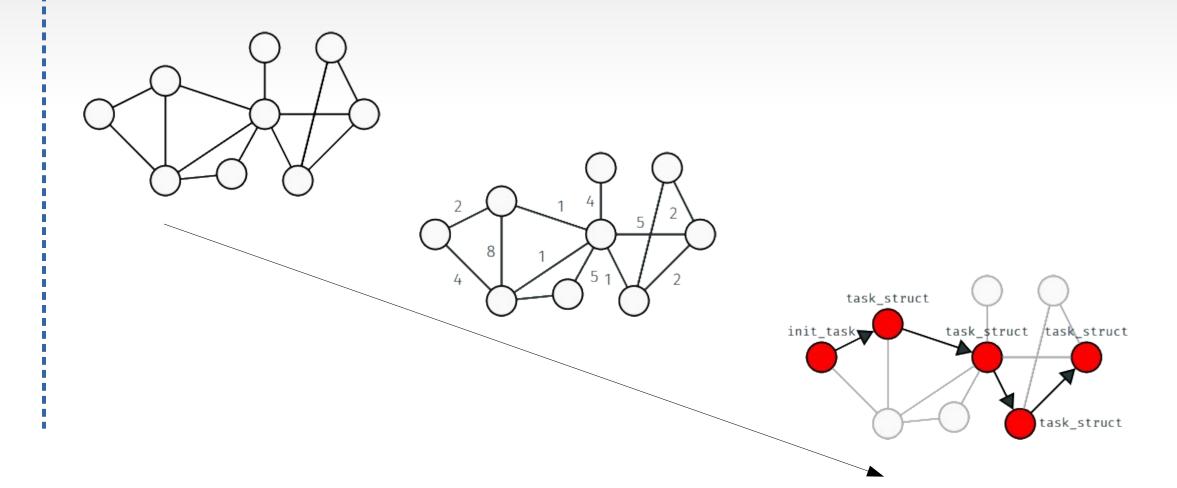
53% of nodes (and 96% of those used by Volatility) are part of a single stronglyconnected component

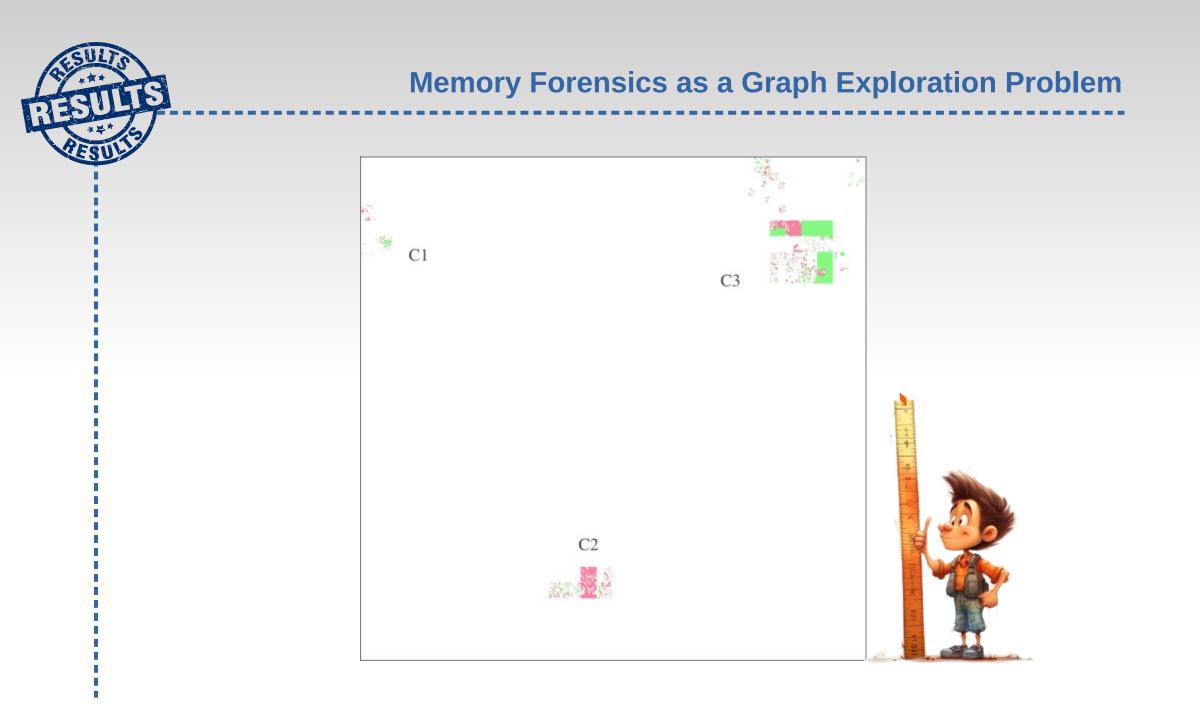




Memory Forensics as a Graph Exploration Problem

Path comparison based on different metrics



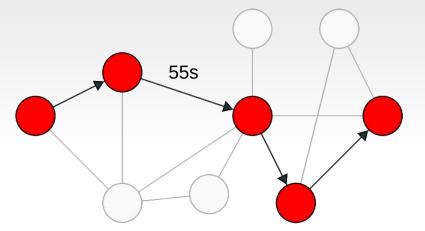


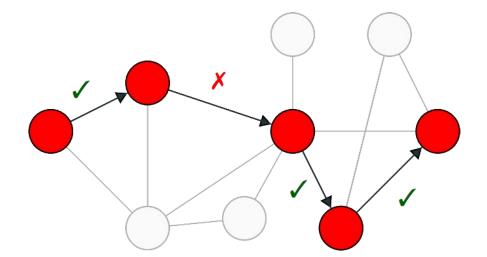


Memory Forensics as a Graph Exploration Problem

Atomicity

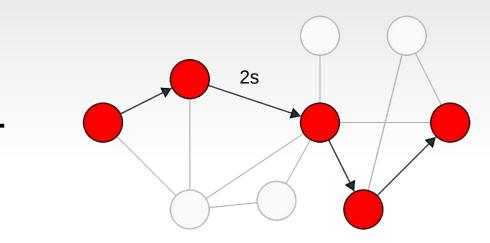
(distance between two structures)





Stability

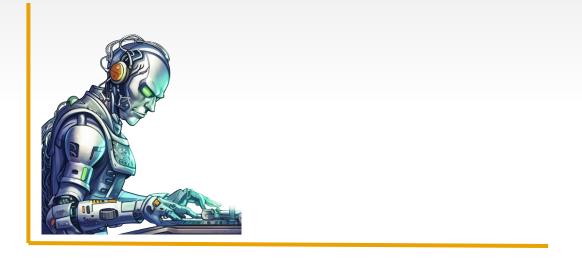
(how often the pointers are updated)





Memory Forensics as a Graph Exploration Problem

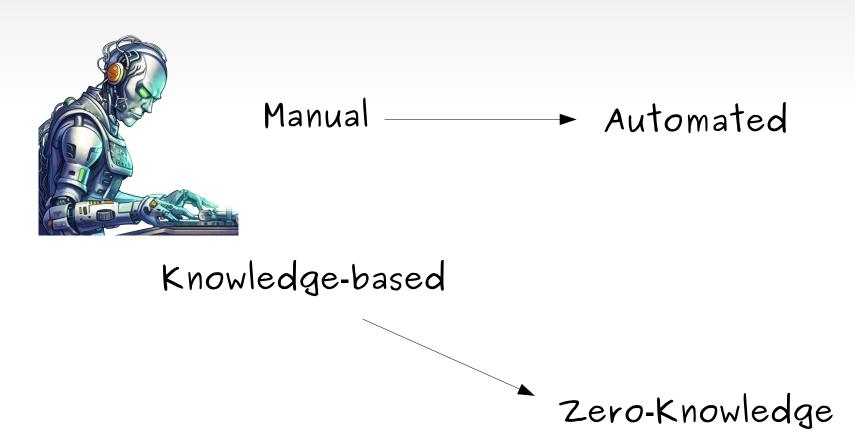
Want to know more? 💋 u s e n i x THE ADVANCED COMPUTING SYSTEMS ASSOCIATION Back to the Whiteboard: a Principled Approach for the Assessment and Design of Memory **Forensic Techniques** Fabio Pagani and Davide Balzarotti, EURECOM https://www.usenix.org/conference/usenixsecurity19/presentation/pagani This paper is included in the Proceedings of the 28th USENIX Security Symposium. August 14–16, 2019 • Santa Clara, CA, USA 978-1-939133-06-9

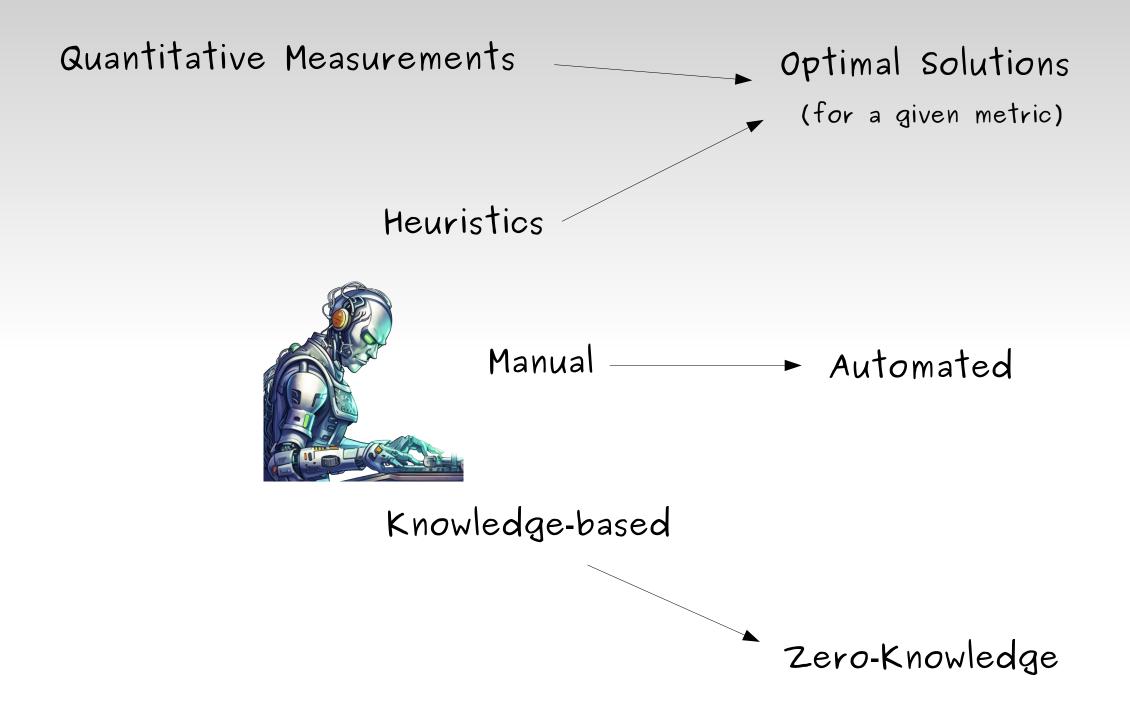


Memory Forensics 2.0 ?



Manual – Automated







"Wet the Appetite" by Midjourney



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http://s3.eurecom.fr/~balzarot